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## HOW TO SUPPLY ASSEMBLY LINES

Great efforts are currently dedicated to cost engineering in industry. The lean production approach is a good way to lower production costs. Material flows and treatment should be planned and controlled neither too early nor too late, because early production leads to an increase of stocks of materials and late production to a backlog of orders. This task is especially important in serial production, e.g. in the automotive industry. A buffer stock could be connected with assembly lines in many ways, such as handcarts, fork lift trucks or special cableways. A modern method of material supply in an assembly workshop is the usage of tow trains consisting of a tow truck and several carts. The number of carts in a train can change but it cannot be high because of turnings and crossings. The task is to find and optimize transport routes. This task is considered to be NP-hard. The paper describes several approaches to solving the task.

**Keywords:** Material supplies, heuristic optimization, assembly lines, milk run VRP problem.

### 1. Task specification

There are two main possibilities for building tow trains in a central store:

- Using static milk run
- Using dynamic milk runs

The first case expects that transport routes are constantly designed. Assembly lines are operated with trains which run on fixed routes, sometimes in fixed periods. Optimization tasks are focused on finding good coverage of an assembly workshop with transport routes, and estimating the optimal number of trains and their lengths. Another task is to find a good solution for the bin packing problem of a cart. The optimization of such a system has been made by Ulrych and Raska [1].

A more interesting way of supplying assembly lines is the creation of routes and lengths of tow trains according to the momentary situation. There are demands from some assembly points for a new delivery of a material batch. The delivery must be in time. The criterion of optimization is a minimum sum of the length of all routes. The task is open, which means that during processing of an optimal schedule of tow trains new demands can appear. Trains on routes must carry out their orders but the rest of the demands are added to with new demands and newly optimized.

An example of such an assembly system is shown in Fig. 1. There is a central store, several assembly lines with work places and supply points. Tow trains can go in lanes. Lanes are mostly one-way. Trains cannot be overtaken and must wait when another train provides a supply point.

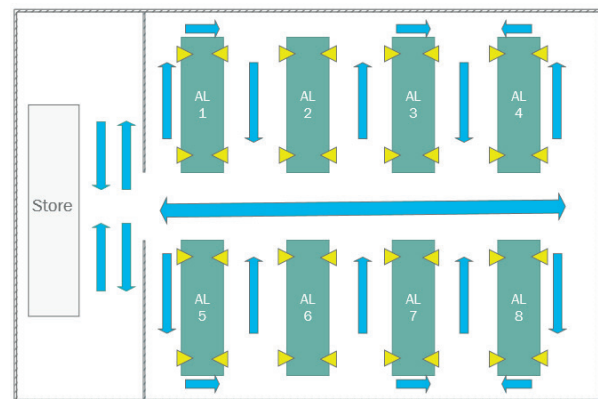


Fig. 1 Model of an assembly line (AL) supply system

This model can be represented in the form of a graph (Fig. 2). In this example, U00 represents the store, vertices Uxx supply points on assembly lines. Edges are evaluated by distances between them. Transport demands of U01, U05, U17, U18, U11, U13 and U14 can be covered by two cycles (red and blue) which represent routes of two supply trains.

If the length of a tow train is  $n$  carts then the task of optimizing the sum of lengths of routes can be represented by covering the graph with cycles going through maximal  $n$  vertices of the graph.

Groups of  $n$  demands ( $1 < n < 20$ ) are randomly generated from vertices needing service by trains. Every group can be considered as a subgraph. The task is to cover the subgraph with cycles of a maximum of 5 vertices (start in the central store and

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**Tow trains in assembly line systems**

Let's say that the length of a train is maximal n. For less or equal demands than n all permutations can be computed (the practical length is under 7 carts). For more than n demands it is possible to randomly select several n-tuples of demands. Each n-tuple is optimized by computing all permutations. The random selection of n-tuples could be repeated many times. The number of repetitions is limited only by computer time (for example under 3 seconds).

**6. Approval of hypotheses**

**Solution by brute force**

The optimization problem is considered to be NP-hard. It is possible to solve it by brute force trying all permutations of high order up to 12 demands. To verify it, a test program has been developed. Results are shown in Table 1.

Times for solutions with brute force Table 1

n	factorial	Sec	min
2	2	0.000	0.000
3	6	0.000	0.000
4	24	0.000	0.000
5	120	0.000	0.000
6	720	0.000	0.000
7	5 040	0.000	0.000
8	40 320	0.003	0.000
9	362 880	0.029	0.000
10	3 628 800	0.277	0.005
11	39 916 800	3.020	0.050
12	479 001 600	36.082	0.601
13	6 227 020 800	489.931	8.166
14	87 178 291 200	7 000.000	116.667

**Stacker**

Simulation programs have been developed to validate the hypotheses.

It is possible to consider that demands for a stacker in a FMS have regular distribution of demands, which means that horizontal positions of all start and end points of demands could appear with the same probability. The task is to minimize the sum of the length of unloaded moves. With no optimization (FIFO - First In First Out), the ratio between unloaded and loaded moves is 100%.

Four strategies of optimization have been tested (Fig. 3):

- No optimization (FIFO).
- Selection of the nearest demand (Proximate).

- Repeated selection of random permutation (Random permutation).
- Brute force (all possible permutation up to 10 demands, then the computing time rises too much).

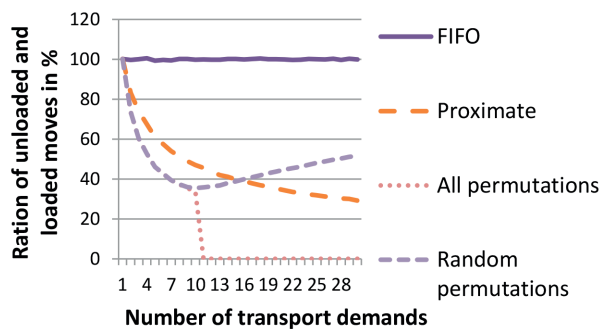


Fig. 3 Optimization of a stacker in a FMS

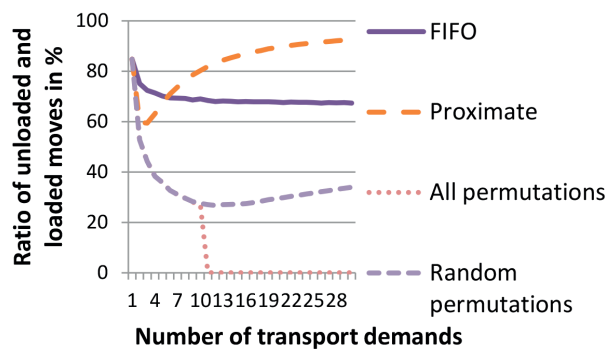


Fig. 4 Optimization of a stacker in a store

Stackers are often used in stores (Fig. 4). Moves of a stacker in a store are asymmetric which means that it either starts at one end (input place of a rack) and ends anywhere in a rack, or starts anywhere in a rack and ends at the same end (output place) where the input place is. This case was also simulated and optimized.

The simulations verified the hypotheses.

**Tow trains in assembly line systems**

The maximum length of a train was considered to be 4 carts and 1 tow truck. The optimum in a 4-tuple was found by testing all 24 possibilities. The task was to divide all transport demands into 4-tuples. Two strategies were used:

- selection of the proximate demand (deterministic),
- repeated random selection of 4-tuples, evaluation of lengths of routes for all tuples, the selection of the best selection (minimum length).

The worst selection was chosen for comparison.

This simulation verified that this method can decrease lengths of routes to about 70% of an unoptimized strategy (Fig. 5).

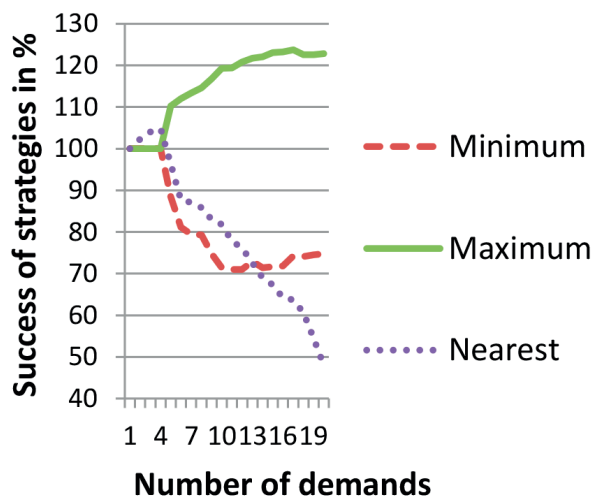


Fig. 5 Success of strategies (nearest demand, min and max random selection)

### 7. Conclusion

The research carried out on an “ad hoc” proposed model proved the hypothesis that a dynamic optimization of train creation and specification of supply routes of trains can be made in real time using relatively simple methods which are not

critically dependent on the accuracy of input data with good optimization results. Means of transportation and manpower can be saved in this way, and peaks of transport demands can be solved more quickly.

The model presented here has been developed theoretically, but inspired by an actual production system model of an assembly line. Subsequent research will be based on the one hand on optimization of larger production systems with the hypothesis of division of the optimized system into two or more parts, and on the other hand on specification of accurate data of an actual production system and experiments leading to optimization of its supply demands.

The optimization of a stacker in a FMS or in a store has been presented for comparison but this strategy was used in several production systems many years ago.

Repeated random selections could be a good solution when optimization tasks are considered for NP-complete. As has been shown above, some “ad hoc” deterministic strategies can also be used and compared with random selections.

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